



CPS II

15-150 M21

Lecture 0630
30 June 2021

0 CPS Predicates

Invariant $p : t \text{ pred}$ iff $p : t \rightarrow \text{bool}$ and p is total

0630.0 (cpsPred.sml)

```
3 type 'a pred = 'a -> bool
```

0630.1 (cpsPred.sml)

```
6 fun simpleDivBy (n:int):int pred =  
7   fn x => (x mod n)=0  
8 val divByTwo0 = simpleDivBy 2  
9 val divByThree0 = simpleDivBy 3  
10 val divByFour0 = simpleDivBy 4  
11 val divByFive0 = simpleDivBy 5  
12 val divBySix0 = simpleDivBy 6
```

0630.2 (cpsPred.sml)

```
16 fun filter (p : 'a pred) ([] : 'a list) = []  
17   | filter p (x::xs) =  
18     if p(x)  
19     then x::filter p xs  
20     else filter p xs
```

```
filterCPS0 : 'a pred -> 'a list -> ('a list -> 'b)  
  -> 'b
```

REQUIRES: true

ENSURES:

$$\text{filterCPS0 } p \ L \ k \cong k(\text{filter } p \ L)$$

0630.3 (cpsPred.sml)

```
24 fun filterCPS0 (p : 'a pred) ([] : 'a list) k
25   = k []
26   | filterCPS0 p (x::xs) k =
27     if p(x)
28     then filterCPS0 p xs (fn l => k(x::l))
29     else filterCPS0 p xs k
```

Crazy idea:

What if the predicate took in a continuation?


```
(* p : int pred (i.e. p:int->bool, p total) *)  
fun p (x:int):bool = ...
```

Accept: $p(x) \implies \text{true}$

Reject: $p(x) \implies \text{false}$

```
(* p :      int  
      -> (unit -> 'a)  
      -> (unit -> 'a)  
      -> 'a                                     *)  
fun p x sc fc = ...
```

Accept: $p\ x\ sc\ fc \implies sc()$

Reject: $p\ x\ sc\ fc \implies fc()$

Invariant $P : (t1, t2)$ cpsPred iff for all x, sc, fc ,

$$P \ x \ sc \ fc \implies sc() \quad \text{or} \quad P \ x \ sc \ fc \implies fc()$$

0630.4 (cpsPred.sml)

```

38 type ('a, 'b) cpsPred =
39   'a -> (unit -> 'b) -> (unit -> 'b) -> 'b

```

0630.5 (cpsPred.sml)

```
43 fun cpsDivBy (n:int) : (int,'a) cpsPred =  
44     fn x => fn sc => fn fc =>  
45         if (x mod n)=0 then sc() else fc()  
46 fun divByTwo1 x = cpsDivBy 2 x  
47 fun divByThree1 x = cpsDivBy 3 x  
48 fun divByFour1 x = cpsDivBy 4 x  
49 fun divByFive1 x = cpsDivBy 5 x  
50 fun divBySix1 x = cpsDivBy 6 x
```

```
filterCPS1 : ('a,'b) cpsPred -> 'a list -> ('a  
list -> 'b) -> 'b
```

REQUIRES: true

ENSURES: `filterCPS1 P L k` evaluates to `k(L')`, where `L'` is the sublist of `L` containing all those `x` such that `P accepts x`.

Live Coding

```
54 fun filterCPS1 (P : ('a,'b) cpsPred)
55           ([] : 'a list)
56           (k : 'a list -> 'b)
57     = k []
58 | filterCPS1 P (x::xs) k =
59     P x
60     (fn () =>
61       filterCPS1
62         P
63         xs
64         (fn res => k(x::res)))
65     (fn () => filterCPS1 P xs k)
```

Idea: When the “success” case happens, instead of just passing in a `unit` to the success continuation, let’s pass in some kind of data, a “witness” to or “evidence” of the success.

```
(* P :      'a
   -> ('e -> 'b)
   -> (unit -> 'b)
   -> 'b                                     *)
fun p x sc fc = ...
```

Accept (with evidence e): $P\ x\ sc\ fc \implies sc\ (e)$

Reject: $P\ x\ sc\ fc \implies fc\ ()$

Invariant $P : (t1, et, t2) \text{ evidPred}$ iff for all x, sc, fc ,
 $P \ x \ sc \ fc \implies sc(e)$ for some $e : et$ or $P \ x \ sc \ fc \implies fc()$

0630.7 (cpsPred.sml)

```
73 type ('a, 'e, 'b) evidPred =  
74   'a -> ('e -> 'b) -> (unit -> 'b) -> 'b
```


0630.8 (cpsPred.sml)

```
77 fun evidDivBy (n:int): (int,int,'a) evidPred =  
78     fn x => fn sc => fn fc =>  
79     if (x mod n)=0 then sc(x div n) else fc()  
80 fun divByTwo2 x = evidDivBy 2 x  
81 fun divByThree2 x = evidDivBy 3 x  
82 fun divByFour2 x = evidDivBy 4 x  
83 fun divByFive2 x = evidDivBy 5 x  
84 fun divBySix2 x = evidDivBy 6 x
```

```
filterCPS2 : ('a, 'e, 'b) evidPred -> 'a list -> (('a * 'e) list -> 'b) -> 'b
```

REQUIRES: true

ENSURES: `filterCPS2 P L k` evaluates to `k(L')`, where `L'` consists of all pairs `(x, e)` where `x` is an element of `L` and `P x sc fc \implies sc(e)`.

Live Coding

0630.9 (cpsPred.sml)

```
88 fun filterCPS2
89     (P : ('a,'e,'b) evidPred) [] k
90     = k []
91 | filterCPS2 P (x::xs) k =
92     P x
93     (fn eL => filterCPS2 P xs
94          (fn res => k((x,eL)::res)))
95     (fn () => filterCPS2 P xs k)
```

Find Sublist

We'll be working with values

$P : 'a \text{ list } \text{pred}$

Goal: given $L : t \text{ list}$ and $p : t \text{ list } \text{pred}$, we want to find some sublist L' of L such that p accepts L' ($P \ L' \cong \text{true}$).

```
findSublist0 : 'a list pred -> 'a list -> ('a list  
-> 'b) -> (unit -> 'b) -> 'b
```

REQUIRES: true

ENSURES: $\text{findSublist0 } p \ L \ sc \ fc$ evaluates to $sc \ L'$ for some sublist L' of L such that p accepts L' . If there is no such L' , then $\text{findSublist0 } p \ L \ sc \ fc \implies fc()$.

Live Coding

0630.10 (sublist.sml)

```
3 fun findSublist0 p [] sc fc =  
4     if p [] then sc [] else fc ()  
5 | findSublist0 p (x::xs) sc fc =  
6     findSublist0  
7     (fn l => p(x::l))  
8     xs  
9     (fn l => sc(x::l))  
10    (fn () => findSublist0 p xs sc fc)
```



```
14 fun findSublist2
15     (P : ('a list, 'a list, 'b) evidPred)
16     ([ ] : 'a list)
17     (sc : 'a list -> 'b)
18     (fc : unit -> 'b)
19     : 'b =
20     P [ ] sc fc
21 | findSublist2 P (x::xs) sc fc =
22     findSublist2
23     (fn l => P (x::l))
24     xs
25     (fn l => sc(x::l))
     (fn () => findSublist2 P xs sc fc)
```

5-minute break

1 CPS Iteration

Many of our list functions can fit into the following description:

Step through the list. For each element, either: (a) throw the element away, (b) combine the element into our ongoing accumulation, (c) stop the process, succeeding with the current element, (d) stop the process with a failure

We're going to abstract this into a very general CPS function.

Module: CPSIteration

result type indicates what to do with an element

aux-library/CPSIterate.sml

```
4 datatype result = Accept
5                 | Keep
6                 | Discard
7                 | Break of string
```

A function `check : t -> result` will “govern” the iteration of a `t list`.

```

For : ('a -> result)
  -> 'a list -> ('a -> 'b -> 'b) -> 'b
  -> ('a -> 'c)
  -> (string -> 'c)
  -> ('b -> 'c)
  -> 'c

```

REQUIRES: check is total, combine x total for all x

ENSURES: (For check L combine base success panic return) iterates through L, applying check to each element. It results in either return(z) (where z is all the elements x such that check x is Keep, combined together with base), or success x for some x in L such that check x is Accept, or panic s for some s such that check x \cong Break s for some x in L.

Live Coding


```

53 fun For (check : 'a -> result)
54     (L : 'a list)
55     (combine : 'a -> 'b -> 'b)
56     (base : 'b)
57     (success : 'a -> 'c)
58     (panic : string -> 'c)
59     (return : 'b -> 'c)
60     : 'c
61 =
62 let
63     fun run ([] : 'a list) (k:'b -> 'c) : 'c =
64         k base
65         | run (x::xs) k =
66             (case (check x) of
67                 Accept => success x
68                 | Keep => run xs (k o (combine x))
69                 | Discard => run xs k
70                 | (Break s)=> panic s)
71 in
72     run L return
73 end

```

Example: Prime divisors

0630.13 (iterate.sml)

```
10 fun div_check m 0 = Break "Divide by zero"
11   | div_check m n =
12     case (m div n, m mod n) of
13         (1, 0) => Accept
14         | (_, 0) => Keep
15         | _     => Discard
16 fun div_success n = SOME [n]
17 fun div_combine x xs = x::xs
```

0630.14 (iterate.sml)

```
20 fun primedivisors m =
21   For (div_check m) primes div_combine
22     (> NONE) SOME
```

0630.15 (iterate.sml)

```
26 fun IGNORE x = raise Fail "Ignored"
27 fun KEEP _ = Keep
28 fun DISCARD _ = Discard
29 fun CASE p x = if p x then Keep else Discard
30 fun OPTCASE p x = case p x of
31     (SOME _) => Keep
32     | NONE => Discard
33 val valOf : 'a option -> 'a = Option.valOf
```


Thank you!